


# Probabilistic Method and Random Graphs

## Lecture 4. Large Deviation Theory & Bins and Balls<sup>1</sup>

Xingwu Liu

Institute of Computing Technology  
Chinese Academy of Sciences, Beijing, China

---

<sup>1</sup>The slides are mainly based on Chapter 5 of *Probability and Computing*. 

Questions, comments, or suggestions?

## Two questions

- Do moments uniquely determine the distribution?
- Why are Chernoff bounds so tight?

## Generating functions

Invented by Abraham de Moivre to compute Fibonacci numbers.

Moment generating functions:  $M_X(t) = \mathbb{E}[e^{tX}]$ .

Unique when bounded or convergent around 0: why?

# Chernoff bound in a big picture

## Fundamental laws of probability theory

**Law of large numbers** (Cardano, Jacob Bernoulli 1713, Poisson 1837): The sample average converges to the expected value.

**Central limit theorem** (Abraham de Moivre 1733, Laplace 1812, Lyapunov 1901, Pólya 1920): The arithmetic mean of independent random variables is approximately normally distributed.

$$\lim_{n \rightarrow \infty} \Pr \left( \sqrt{n} \left( \left( \frac{1}{n} \sum_{i=1}^n X_i \right) - \mu \right) \leq x \right) = \Phi \left( \frac{x}{\sigma} \right)$$

## Marvelous but ...

Say nothing about the rate of convergence

## Large deviation theory

How fast does it converge? Beyond central limit theorem

# A glance at large deviation theory

## Motivation

$X_n$ : the number of heads in  $n$  flips of a fair coin.

By the central limit theorem,  $\Pr(X_n \geq \frac{n}{2} + \sqrt{n}) \rightarrow 1 - \Phi(1)$ .

What about  $\Pr(X_n \geq \frac{n}{2} + \frac{n}{3})$ ? Nothing but converging to 0.

## Chernoff bounds say...

$$\Pr(X_n \geq \frac{n}{2} + \frac{n}{3}) \leq \left( \frac{e^{\frac{2}{3}}}{(\frac{5}{3})^{\frac{2}{3}}} \right)^{\frac{n}{2}} \approx e^{-0.092n}.$$

## Actually

Direct calculation shows that

$$\Pr(X_n \geq \frac{n}{2} + \frac{n}{3}) \approx e^{-0.2426n + o(n)} \ll \text{Chernoff bound}.$$

Oh, no!

# Mission of Large Deviation Theory

Find the asymptotic probabilities of *rare events* - how do they decay to 0 as  $n \rightarrow \infty$ ?

*Rare events* mean large deviation.  
So large that CLT is almost useless (deviation up to  $\sqrt{n}$ ).

## Intuition

Inspired by Chernoff bounds, conjecture that probabilities of rare events will be exponentially small in  $n$ :  $e^{-cn}$  for some  $c$ .

Q: Does  $\lim_{n \rightarrow \infty} \frac{1}{n} \ln \Pr(\mathcal{E}_n^{\text{rare}})$  exist? If so, what's it?

# Large Deviation Principle

Simple form (By courtesy of Cramer, 1938)

Let  $X_1, \dots, X_n, \dots \in \mathbb{R}$  be i.i.d. r.v. which satisfy  $\mathbb{E}[e^{tX_1}] < \infty$  for  $t \in \mathbb{R}$ . Then for any  $t > \mathbb{E}[X_1]$ , we have

$$\lim_{n \rightarrow \infty} \frac{1}{n} \ln \Pr\left(\sum_{i=1}^n X_i \geq tn\right) = -I(t),$$

where

$$I(t) \triangleq \sup_{\lambda > 0} \lambda t - \ln \mathbb{E}[e^{\lambda X_1}].$$

Remark

$I(\cdot)$ : rate function.

Many variants: the factor  $\frac{1}{n}$ ,  $tn$  in the events, random variables

# Large Deviation Principle: Proof

## Large Deviation Principle

$$\lim_{n \rightarrow \infty} \frac{1}{n} \ln \Pr(\sum_{i=1}^n X_i \geq tn) = -(\sup_{\lambda > 0} \lambda t - \ln \mathbb{E}[e^{\lambda X_1}]).$$

## Proof: Upper bound

Let  $Y_n = \frac{\sum_{i=1}^n X_i}{n}$ ,  $M(\lambda) = \mathbb{E}[e^{\lambda X_1}]$ , and  $\psi(\lambda) = \ln M(\lambda)$ .

$$\Pr(Y_n \geq t) \leq e^{-\lambda nt} (M(\lambda))^n \text{ for any } \lambda \geq 0.$$

$$\frac{1}{n} \ln \Pr(Y_n \geq t) \leq -\lambda t + \psi(\lambda).$$

$$\frac{1}{n} \ln \Pr(Y_n \geq t) \leq -\sup_{\lambda \geq 0} (\lambda t - \psi(\lambda)).$$



# Large Deviation Principle: Proof

## Lower bound

The maximizer  $\lambda_0$  of  $\lambda t - \psi(\lambda)$  satisfies  $t = \int \frac{x e^{\lambda_0 x}}{M(\lambda_0)} d\mu(x)$ .

Let  $d\mu_0(x) = \frac{e^{\lambda_0 x}}{M(\lambda_0)} d\mu(x)$ . Its expectation  $\int x d\mu_0(x) = t$ .

Let  $A = \{Y_n \geq t\} \subseteq \mathbb{R}^n$ ,  $A_\delta = \{Y_n \in [t, t + \delta]\} \subseteq \mathbb{R}^n$ .

$$\begin{aligned} \Pr_\mu(A) &\geq \Pr_\mu(A_\delta) = \int_{A_\delta} \prod_{i=1}^n d\mu(x_i) \\ &= \int_{A_\delta} (M(\lambda_0))^n e^{-\lambda_0 \sum_{i=1}^n x_i} \prod_{i=1}^n d\mu_0(x_i) \\ &\geq (M(\lambda_0) e^{-\lambda_0(t+\delta)})^n \Pr_{\mu_0}(A_\delta). \end{aligned}$$

Applying CLT to  $\mu_0$ , we have  $\lim_{n \rightarrow \infty} \Pr_{\mu_0}(A_\delta) = \frac{1}{2}$ .

$\lim_{n \rightarrow \infty} \frac{1}{n} \ln \Pr(Y_n \geq t) \geq \psi(\lambda_0) - (t + \delta)\lambda_0$ , and let  $\delta \rightarrow 0$ .

## Large deviation theory vs CLT

Seemingly easy to get exponential decay in many cases, but hard to calculate.

Chernoff bounds concern large deviation

- Con: Generally weaker
- Pro: Always holds, not just asymptotically

Key assumption

**Independence!**

## Main idea

Approximation with independence.

## Focus

Approximation.

# The Bins-and-Balls Model

General setting:  $(m, n)$ -model

## Extension

Multiple choice, limited capacity of bins ...

## Applications

**Load balancing:** balls = jobs, bins = servers;

**Data storage:** balls = files, bins = disks;

**Hashing:** balls = data keys, bins = hash table slots;

**Coupon Collector:** balls = coupons; bins = coupon types.

Number of balls in any bin:  $\text{Bin}(m, \frac{1}{n})$ .

Numbers of balls in multiple bins: not independent. Why?

Application: time complexity of bucket-sort

**Bucket-sort:** Given  $n = 2^m$  integers from  $[0, 2^k)$  with  $k > m$ , first allocate the integers to  $n$  bins, followed by sorting each bin.

**Expected time complexity:**  $n + \mathbb{E}[\sum_{i=1}^n X_i^2] = n + n\mathbb{E}[X_1^2]$ .  
 $X_1 \sim \text{Bin}(n, \frac{1}{n})$ , so  $\mathbb{E}[X_1^2] = 2 - \frac{1}{n}$ .

# Topics of Bins-and-Balls Model

## The distribution of

Number of balls in a certain bin

Maximum load

Number of bins containing  $r$  balls

...

## Max. load: when does it exceed 1 w.h.p.?

The probability that max. load is 1 is

$$\left(1 - \frac{1}{n}\right)\left(1 - \frac{2}{n}\right) \cdots \left(1 - \frac{m-1}{n}\right) \leq \prod_{i=1}^{m-1} e^{-\frac{i}{n}} \approx e^{-\frac{m^2}{2n}}.$$

It is less than  $\frac{1}{2}$  if  $m \geq \sqrt{2n \ln 2}$

## Birthday paradox

$n = 365, m \geq 22.49$

# Max load: $(n, n)$ -model

Asymptotically,  $\Pr(L \geq 3 \frac{\ln n}{\ln \ln n}) \leq \frac{1}{n}$

## Proof

$X_i$ : the number of balls in bin  $i$ .

$$\Pr(X_1 \geq k) \leq \binom{n}{k} \frac{1}{n^k} \leq \frac{1}{k!}.$$

$$\frac{k^k}{k!} < \sum_i \frac{k^i}{i!} = e^k \Rightarrow \frac{1}{k!} \leq \left(\frac{e}{k}\right)^k.$$

$$\begin{aligned} \Pr\left(L \geq 3 \frac{\ln n}{\ln \ln n}\right) &\leq n \left(\frac{e \ln \ln n}{3 \ln n}\right)^{3 \frac{\ln n}{\ln \ln n}} \\ &\leq n \left(\frac{\ln \ln n}{\ln n}\right)^{3 \frac{\ln n}{\ln \ln n}} \\ &\leq e^{\ln n + (\ln \ln \ln n - \ln \ln n) \frac{3 \ln n}{\ln \ln n}} \leq \frac{1}{n}. \end{aligned}$$

# Number of bins having load $r$ : $(m, n) - model$

$r = 0$

The distribution of  $X_i$ 's are identical:  $Bin(m, \frac{1}{n})$ .

$$\Pr(X_i = 0) = \left(1 - \frac{1}{n}\right)^m \approx e^{-\frac{m}{n}}.$$

Expected number of empty bins is about  $ne^{-\frac{m}{n}}$ .

Load= $r$

$$\Pr(X_i = r) = \binom{m}{r} \frac{1}{n^r} \left(1 - \frac{1}{n}\right)^{m-r}.$$

When  $r \ll \min\{m, n\}$ ,  $\Pr(X_i = r) \approx e^{-\frac{m}{n}} \frac{\left(\frac{m}{n}\right)^r}{r!}$ .

Expected number of load- $r$  bins is about  $ne^{-\frac{m}{n}} \frac{\left(\frac{m}{n}\right)^r}{r!}$ .

Poisson distribution

$$\sum_j e^{-\mu} \frac{\mu^j}{j!} = 1 \text{ due to } e^x = \sum_j \frac{x^j}{j!}.$$

Nonnegative-integer-valued r.v.  $X_\mu$ :  $\Pr(X_\mu = j) = e^{-\mu} \frac{\mu^j}{j!}$ .



# Basic Properties of Poisson distribution

## Low-order moments

$$\mathbb{E}[X_\mu] = \text{Var}[X_\mu] = \mu.$$

## Moment generation function

$$M_{X_\mu}(t) = \mathbb{E}[e^{tX_\mu}] = \sum_k \frac{e^{-\mu} \mu^k}{k!} e^{tk} = e^{\mu(e^t-1)}.$$

## Additive

By uniqueness of moment generation functions,  
 $X_{\mu_1} + X_{\mu_2} = X_{\mu_1+\mu_2}$  if independent.

## Chernoff-like bounds

1. If  $x > \mu$ , then  $\Pr(X_\mu \geq x) \leq \frac{e^{-\mu}(e\mu)^x}{x^x}$ .
2. If  $x < \mu$ , then  $\Pr(X_\mu \leq x) \leq \frac{e^{-\mu}(e\mu)^x}{x^x}$ .

## Occurrences of **rare events** during a fixed interval

- Typos per page in printed books.
- Number of bomb hits per  $0.25\text{km}^2$  in South London during World War II.
- The number of goals in sports involving two competing teams.
- *The number of soldiers killed by horse-kicks each year in Prussian cavalry corps in the (late) 19th century.*

## Story of Poisson distribution

1837, Poisson, *Research on the Probability of Judgments in Criminal and Civil Matters*.

Appeared in 1711, de Moivre. (Stigler's law of eponymy, 1980)

First practical application (next page)

# First practical application of Poisson distribution

## Reliability engineering

- Ladislaus Bortkiewicz (1868-1931)
  - Russian economist and statistician of Polish ancestry, mostly lived in Germany
  - Famous for Poisson distribution and Marxian economics
- The book *The Law of Small Numbers*, 1898
- Annual Horse-kick data of 14 cavalry corps over 20 years
- Events with low probability in a large population follow a Poisson distribution

No. deaths $k$	Freq.	Poisson approx. $200 \times \mathbb{P}(\text{Poi}(0.61) = k)$
0	109	108.67
1	65	66.29
2	22	20.22
3	3	4.11
4	1	0.63
5	0	0.08
6	0	0.01

① `http://willperkins.org/6221/slides/largedeviations.pdf`