# Probabilistic Method and Random Graphs 

## Lecture 12. A Brief Introduction to Markov Chains ${ }^{1}$

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[^0] Chains And Random Walks by Takis Konstantopoulos.

## Preface

Questions, comments, or suggestions?

## A recap of Lovász local lemma

## Mission

- Do events $A_{1}, \ldots A_{n}$ satisfy $\operatorname{Pr}\left(\bigcup_{i=1}^{n} A_{i}\right)<1$ ?

Symmetric version: $\operatorname{Pr}\left(\bigcup_{i=1}^{n} A_{i}\right)<1$ when

- edp $\leq 1$ for all $i$, with $p=\max _{i} \operatorname{Pr}\left(A_{i}\right), d=\max _{i}\left|\Gamma\left(A_{i}\right)\right|$

Asymmetric version: $\operatorname{Pr}\left(\bigcup_{i=1}^{n} A_{i}\right)<1$ when

- $\forall i, \sum_{A_{j} \in \Gamma\left(A_{i}\right)} \operatorname{Pr}\left(A_{j}\right) \leq \frac{1}{4}$, or
- $\exists x_{1}, \ldots x_{n} \in(0,1)$ s.t. $\forall i, \operatorname{Pr}\left(A_{i}\right) \leq x_{i} \prod_{A_{j} \in \Gamma\left(A_{i}\right)}\left(1-x_{j}\right)$
- Shearer's bound is tight
- Moser-Tardos algorithm is efficient up to Shearer's bound


## An overall review of probabilistic method

Handling dependence, exploiting independence

- Counting (union bound): mutually exclusive
- First moment: linearity doesnt care dependence
- Second moment: pairwise dependence
- LLL: global dependence

Continue this trend in stochastic process

## Markov Chains

## Informal definition

A mathematical model of a random phenomenon evolving with time such that the past affects the future only through the present

Time can be discrete or continuous (Markov process)

## Debut of the concept of Markov chains

Andrey Markov. Extension of the law of large numbers to dependent quantities, Izvestiia Fiz.-Matem. Obsch. Kazan Univ., (2nd Ser.), 15(1906), pp. 135-156

From an individual to a sequence of random variables

- Asymptotical behavior matters


## Andrey Andreyevich



Example: a mouse in cage

## 1 <br>  <br> 2

Behavior of the mouse (transition diagram): $\alpha=0.05, \beta=0.99$


## Example: a mouse in cage

Behavior of the mouse (transition matrix)

$$
\mathbf{P}=\left(\begin{array}{cc}
1-\alpha & \alpha \\
\beta & 1-\beta
\end{array}\right)=\left(\begin{array}{ll}
0.95 & 0.05 \\
0.99 & 0.01
\end{array}\right)
$$

## Interesting questions

- How long does it take for the mouse, on the average, to move from cell 1 to cell 2?
- Easy to solve due to the geometric distribution
- How often is the mouse in room 1?
- Hard to answer it in one minute


## Example: insurance company's puzzle

## Human health on a monthly basis



## Transition matrix

$P=\left(\begin{array}{ccc}0.69 & 0.3 & 0.01 \\ 0.8 & 0.1 & 0.1 \\ 0 & 0 & 1\end{array}\right)$

What is the distribution of the lifetime of a currently healthy one?

## Formal definition of Markov Chains

## General setting

- A sequence of random variables $\left\{X_{n}: n \in \mathbb{N}\right\}$
- For all $n, X_{n}$ is defined on the same state space $S$
- Any $s \in S$ is called a state

> Markov property
> $\operatorname{Pr}\left(X_{n+1}=x_{n+1} \mid X_{n}=x_{n}, \ldots X_{0}=x_{0}\right)=\operatorname{Pr}\left(X_{n+1}=x_{n+1} \mid X_{n}=\right.$ $\left.x_{n}\right)$, for any $n \in \mathbb{N}$ and $x_{0}, \ldots x_{n} \in S$

The future is independent of the past, given the present state

## Homogeneous

$\operatorname{Pr}\left(X_{n+1}=y \mid X_{n}=x\right)$ is independent of $n$, denoted by $p_{x y}$

Focus on homogeneous Markov chains

## Representation of a Markov chain

## Transition diagram

Weighted directed graph $G=(V, E, W)$

- $V=S$, the state space
- $e_{i j} \in E$ if and only if $p_{i j} \triangleq \operatorname{Pr}\left(X_{t}=j \mid X_{t-1}=i\right)>0$
- $W: e_{i j} \mapsto p_{i j}$


## This provides intuition

- Example: state reachability is reachability over the graph


## Transition matrix

$P=\left(p_{i j}\right)_{i, j \in S}$, all entries are nonnegative, $\sum_{j} p_{i j}=1$
This enables calculation

- Example: $P^{(n)}=P^{n}$, where $P^{(n)}=\left(p_{i j}^{(n)}\right)_{i, j \in S}$,

$$
p_{i j}^{(n)} \triangleq \operatorname{Pr}\left(X_{n}=j \mid X_{0}=i\right)
$$

## Multistep transition matrix

$$
P^{(n)}=P^{n}
$$

Proof by induction on $n$.
Remark: a summand of $p_{i j}^{(n)}$ corresponds to a path from $i$ to $j$ whose length is $n$

State distribution at time $t$
Given initial distribution $\pi, \pi^{(t)}=\pi P^{(t)}=\pi P^{t}$

## Interesting questions

- Can a state $j$ be reached from $i$ ?
- If yes, when?
- What's the state distribution at any $t$ ?
- What's the distribution in the long run (average frequency)?


## Reachability

Equivalent conditions of reaching $j$ from $i$

- There is a directed path in $G$ from $i$ to $j$
- $p_{i j}^{(n)}>0$ for some $n$

Denoted by $i \rightsquigarrow j$

## Communicating states

$i \rightsquigarrow j$ if $i \rightsquigarrow j$ and $j \rightsquigarrow i$
Communicating classes: equivalence classes of $\longleftrightarrow \rightsquigarrow$

- Strongly connected components of $G$


## Period

## The life style of a pig


$p_{i i}^{(n)}>0$ only if $n$ is even. It is periodic

## The period of state $i$ of a Markov chain

$d_{i}$ is the GCD of $D_{i} \triangleq\left\{n \geq 1: p_{i i}^{(n)}>0\right\}$.
If $d_{i}=1, i$ is said to be aperiodic

## Communicating states have the same period

## Theorem

If $i \nrightarrow j$, then $d_{i}=d_{j}$

## Proof

- Since $i \stackrel{\text { ↔ }}{ } j, p_{i j}^{(s)}>0$ and $p_{j i}^{(t)}>0$ for some $s, t>0$
- $p_{i i}^{(s+t)} \geq p_{i j}^{(s)} p_{j i}^{(t)}>0$, so $d_{i}$ divides $s+t$
- For any $n \in D_{j}, p_{i i}^{(s+n+t)} \geq p_{i j}^{(s)} p_{j j}^{(n)} p_{j i}^{(t)}>0$, so $d_{i}$ divides $s+n+t$
- Since $d_{i}$ divides $s+t, d_{i}$ divides $n$
- $d_{i}$ divides $d_{j}$
- Symmetrically, $d_{j}$ divides $d_{i}$
- $d_{j}=d_{i}$


## Theorem

If $i$ is aperiodic, $p_{i i}^{(n)}>0$ for all large enough $n$

## Proof

- Choose $n_{1}, n_{2} \in D_{i}$ s.t. $n_{2}-n_{1}=1$
- For any $n$, there are integers $q$ and $r<n_{1}$ s.t. $n=q n_{1}+r$
- $n=q n_{1}+r\left(n_{2}-n_{1}\right)=(q-r) n_{1}+r n_{2}$
- When $n$ is large enough, $q-r>0$
- $p_{i i}^{(n)} \geq\left(p_{i i}^{\left(n_{1}\right)}\right)^{q-r}\left(p_{i i}^{\left(n_{2}\right)}\right)^{r}>0$


## Hitting time

## Definition

$T_{i j}$ : the first time that $j$ is reached when the initial state is $i$

- $f_{i j}^{(n)} \triangleq \operatorname{Pr}\left(T_{i j}=n\right)=\operatorname{Pr}\left(X_{n}=j, X_{k} \neq j, 1 \leq k<n \mid X_{0}=i\right)$
- $f_{i j} \triangleq \sum_{n} f_{i j}^{(n)}$


## Recurrency

If $f_{i i}=1$, the state $i$ is recurrent (otherwise, transient)

- Furthermore, if $\mathbb{E}\left[T_{i i}\right]<\infty, i$ is positive recurrent
- Otherwise, it is null recurrent


## Example

Human health chain, pig life style chain, and more

## Decision theorem of recurrency

The following conditions are equivalent
(1) $i$ is recurrent
(2) $\sum_{n} p_{i i}^{(n)}=\infty$
(3) $\mathbb{E}\left[J_{i} \mid X_{0}=i\right]=\infty, J_{i}$ is the number of times $i$ is reached
(9) $\operatorname{Pr}\left(J_{i}=\infty \mid X_{0}=i\right)=1$

Proof: $2 \Leftrightarrow 3$

$$
\begin{aligned}
J_{i} & =\sum_{n} \mathbf{1}\left(X_{n}=i\right) \\
\mathbb{E}\left[J_{i} \mid X_{0}=i\right] & =\mathbb{E}\left[\sum_{n} \mathbf{1}\left(X_{n}=i\right) \mid X_{0}=i\right] \\
& =\sum_{n}^{n} \operatorname{Pr}\left(X_{n}=i \mid X_{0}=i\right) \\
& =\sum_{n} p_{i i}^{(n)}
\end{aligned}
$$

## Proof (continued)

## $1 \Rightarrow 4$

- Let $J_{i}^{(l)}$ be the times of reaching $i$ no earlier than step $l$
- Property: $J_{i}=J_{i}^{(1)}$
- $g_{i i} \triangleq \operatorname{Pr}\left(J_{i}=\infty \mid X_{0}=i\right)=\lim _{k} \operatorname{Pr}\left(J_{i}^{(1)} \geq k \mid X_{0}=i\right)$
- $\left(J_{i}^{(1)} \geq k+1 \mid X_{0}=i\right)=\cup_{l}\left(T_{i i}=l, J_{i}^{(l+1)} \geq k \mid X_{0}=i\right)$
- $\operatorname{Pr}\left(J_{i}^{(1)} \geq k+1 \mid X_{0}=i\right)=f_{i i} \operatorname{Pr}\left(J_{i}^{(1)} \geq k \mid X_{0}=i\right)=f_{i i}^{k+1}$
- $g_{i i}=\lim _{k} f_{i i}^{k}=1$ since $i$ is recurrent


## $4 \Rightarrow 3$

Trivial

## Proof: $2 \Rightarrow 1$

- Chapman-Kolmogorov equation:

$$
p_{i j}^{(n)}=\sum_{k=1}^{n} f_{i j}^{(k)} p_{j j}^{(n-k)}, p_{i i}^{(0)}=1
$$

- For any $N$,

$$
\begin{aligned}
\sum_{n=1}^{N} p_{i i}^{(n)} & =\sum_{n=1}^{N} \sum_{k=1}^{n} f_{i i}^{(k)} p_{i i}^{(n-k)} \\
& =\sum_{k=1}^{N} f_{i i}^{(k)} \sum_{n=k}^{N} p_{i i}^{(n-k)} \\
& =\sum_{k=1}^{N} f_{i i}^{(k)} \sum_{n=0}^{N-k} p_{i i}^{(n)} \\
& \leq \sum_{k=1}^{N} f_{i i}^{(k)} \sum_{n=0}^{N} p_{i i}^{(n)}
\end{aligned}
$$

- $\frac{\sum_{n=1}^{N} p_{i i}^{(n)}}{1+\sum_{n=1}^{N} p_{i i}^{(n)}}=\frac{\sum_{n=1}^{N} p_{i i}^{(n)}}{\sum_{n=0}^{N} p_{i i}^{(n)}} \leq \sum_{k=1}^{N} f_{i i}^{(k)} \leq f_{i i} \leq 1$
- Since $\sum_{n=1}^{N} p_{i i}^{(n)}=\infty$, the lefthand side $\rightarrow 1$ as $N \rightarrow \infty$
- $f_{i i}=1$, so $i$ is recurrent


## Recurrency is preserved by communicating relation

## Theorem

If $i \nrightarrow j$ and $i$ is recurrent, then so is $j$

## Prove

It immediately follows from the above theorem

## A necessary condition of transient states

## Theorem

If $j$ is a transient, $\sum_{n=1}^{\infty} p_{i j}^{(n)}<\infty$ for any $i$

## Proof

- $p_{i j}^{(n)}=\sum_{k=1}^{n} f_{i j}^{(k)} p_{j j}^{(n-k)}, p_{i i}^{(0)}=1$
- For any $N$,

$$
\begin{aligned}
\sum_{n=1}^{N} p_{i j}^{n} & =\sum_{n=1}^{N} \sum_{k=1}^{n} f_{i j}^{(k)} p_{j j}^{(n-k)} \\
& =\sum_{k=1}^{N} \sum_{n=k}^{N} f_{i j}^{(k)} p_{j j}^{(n-k)} \\
& =\sum_{k=1}^{N} f_{i j}^{(k)} \sum_{n=0}^{N-k} p_{j j}^{(n)} \\
& \leq \sum_{k=1}^{N} f_{i j}^{(k)} \sum_{n=0}^{N} p_{j j}^{(n)} \\
\bullet \sum_{n=1}^{N} p_{i j}^{(n)} \leq & \sum_{n=0}^{N} p_{j j}^{(n)} \leq 1+\sum_{n=1}^{N} p_{j j}^{(n)}<\infty
\end{aligned}
$$

## Positive recurrency

Any rule for deciding if a state is positive recurrent?

How to compute the expected hitting time of a positive recurrent state?

## Reference

- Introductory Lecture Notes on Markov Chains And Random Walks by Takis Konstantopoulos
- Baidu Wenku


[^0]:    ${ }^{1}$ The slides are mainly based on Introductory Lecture Notes on Markov

