Probabilistic Method and Random Graphs

Lecture 11a. Lovász Local Lemma

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[1]The slides are mainly based on Chapter 6 of Probability and Computing.

Comments, questions, or suggestions?

Recap of Lecture 10

- Chebyshev's Ine.: $\Pr(|X - \mathbb{E}[X]| \geq a) \leq \frac{\mathrm{Var}[X]}{a^2}$
- $\Pr(X = 0) \leq \Pr(|X - \mathbb{E}[X]| \geq \mathbb{E}[X]) \leq \frac{Var[X]}{\mathbb{E}[X^2]} \leq \frac{\mathrm{Var}[X]}{(\mathbb{E}[X])^2}$
  - When $X \geq 0$, $\Pr(X > 0) > \frac{(\mathbb{E}[X])^2}{\mathbb{E}[X^2]}$
  - $\Pr(X > \theta\mathbb{E}[X]) \geq \frac{(1-\theta)^2(\mathbb{E}[X])^2}{\mathrm{Var}[X]+(1-\theta)^2(\mathbb{E}[X])^2}$

    $\geq (1-\theta)^2 \frac{(\mathbb{E}[X])^2}{\mathbb{E}[X^2]}, \theta \in (0,1)$

Second moment method

Recap: applications

- Distinct Subset sum problem
  - Given $n \in \mathbb{Z}^+$, $f(n)$ is the max k such that:

    $\exists S \subseteq [n], |S| = k, S$ has distinct subset sums
  - Erdős promised 500$: $f(n) = \lfloor \ln_2 n \rfloor + \Theta(1)$
  - $f(n) \le \ln_2 n + \frac{1}{2}\ln_2 \ln_2 n + O(1)$
    - Improve by Hoeffding's inequality?
- Threshold function of $G_{n,p}$
  - Clique number $c(G) \ge 4$
  - $t(n) = n^{-\frac{2}{3}}$
    - The paradigm in estimating the variance of number of 4-cliques

Recap: Lovász local lemma

- Can we avoid all bad events?
- Given bad events $A_1, A_2, \ldots A_n$, is $\Pr(\cap_i \overline{A_i}) > 0$?
  - Applicable to SAT, coloring, Ramsey theory…
  - Challenging when dependent and highly probable
    - Dependency graph (not unique)
- Yes if $\forall i, \Pr(A_i) \leq p, |\Gamma(A_i)| \leq d$ and $4pd \leq 1$
- Any (k,s)-CNF is satisfiable if $s \leq \frac{1}{4}\frac{2^k}{k}$
- $R(k) \geq k 2^{\frac{k}{2}}\left[\frac{\sqrt{2}}{e} + o(1)\right]$. How about $R(k,t)$?

Non-symmetric LLL

- **Theorem**: $\Pr(\cap_i \overline{A_i}) > 0$ if $\forall i, \sum_{j\in\Gamma(A_i)} \Pr\left(A_j\right) < \frac{1}{4}$
  - [Spencer, 1975]
  - The sense of "local"
- Follow the proof of symmetric LLL, with induction on m to show that $\Pr(A|B_1, \ldots, B_m) < 2\Pr(A)$
- Application to $R(k,t)$:

$$R(k,t) > t^{\frac{\binom{k}{2}-2}{k-2}+o(1)}$$ with k fixed and $t \to \infty$

Proof: $R(k,t) > t^{\frac{\binom{k}{2}-2}{k-2}+o(1)}$

- Randomly color edges of K_n, p in red, $(1-p)$ in blue
- S: a k-set of the vertices; T: a t-set of the vertices
- A_S: S is a red clique; B_T: T is a blue clique
- $\Pr(A_S) = p^{\binom{k}{2}}$,$\Pr(B_T) = (1-p)^{\binom{t}{2}}$
- Any event has at most $\binom{t}{2}\binom{n}{k-2}$ neighbors being A_S,

  at most $\binom{n}{t}$ neighbors being B_T
- Let $p = n^{-\epsilon-\beta+\delta}$, $t = n^{\beta+\epsilon}$, $\beta = \frac{k-2}{\binom{k}{2}-2}$, $0 < \delta < \epsilon$,

  we have $\binom{t}{2}\binom{n}{k-2}p^{\binom{k}{2}} + \binom{n}{t}(1-p)^{\binom{t}{2}} < \frac{1}{4}$

Open problem

A stronger non-symmetric LLL

- $\Pr(\cap \overline{A_i}) > 0$ if there are $x_1, x_2, \dots, x_n \in (0,1)$ s.t.
  $\forall i, \Pr(A_i) \leq x_i \prod_{j \in \Gamma(A_i)} (1 - x_j)$
  - Similar proof, but
    - Prove $\Pr\left(A_i \middle| \cap_{j=1}^{t} \overline{B}_j\right) \leq x_i$
    - Use chain rule to lower-bound the numerator $\Pr\left(\cap \overline{C}_j \mid \cap \overline{D_j}\right)$ by $\prod_{j \in \Gamma(A_i)} (1 - x_j)$
  - Spencer, 1977

$$R(k,t) \geq c\left(\frac{t}{\ln t}\right)^{\frac{k+1}{2}}(1-o(1))$$

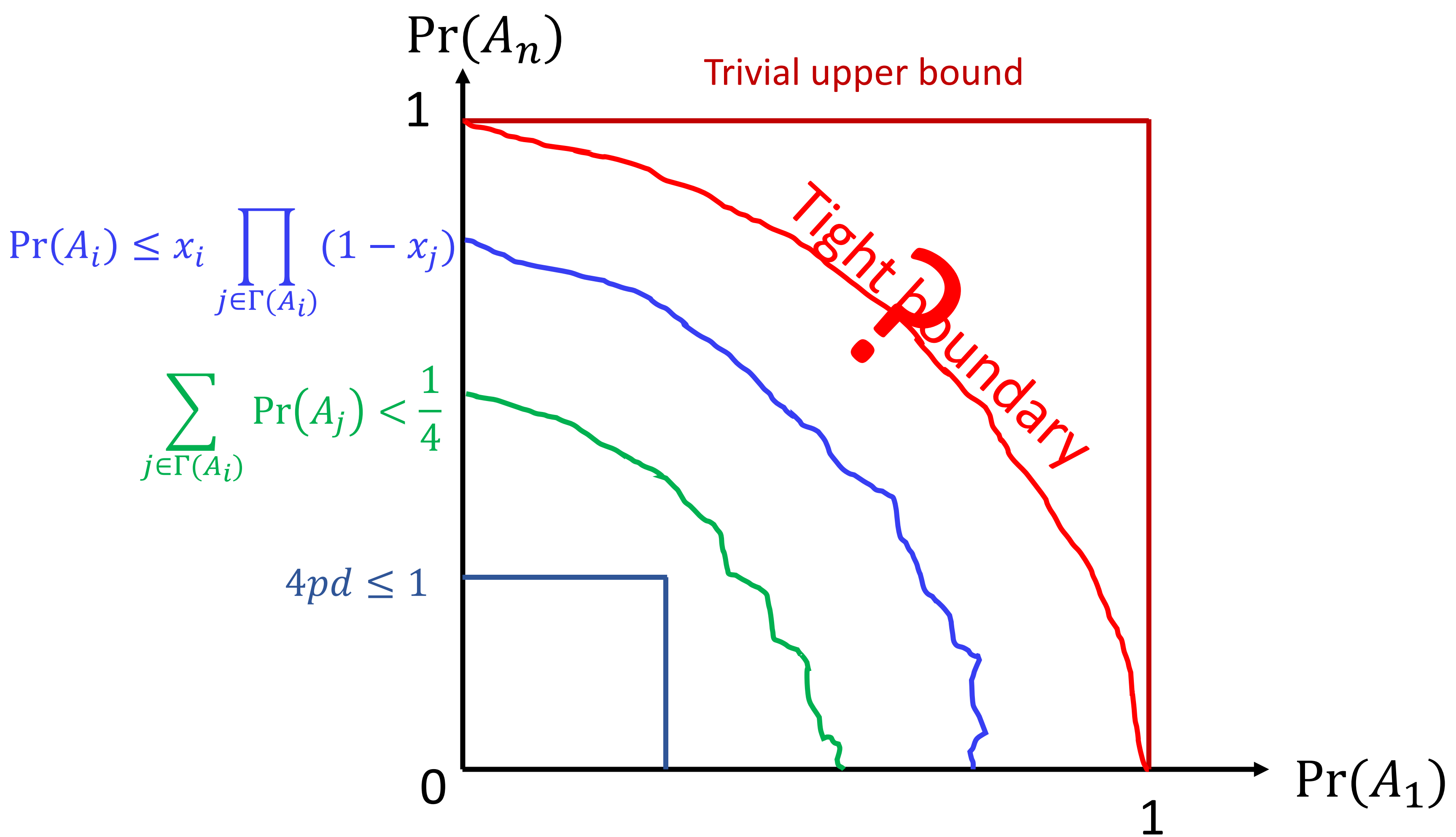
- Follow the proof of $R(k,t) > t^{\frac{\binom{k}{2}-2}{k-2}+o(1)}$
  - Define events A_S and B_T for any k-set S and t-set T
  - Let $p = c_1 n^{-\beta}$, $t = c_2 n^{\beta} \ln n$, $x_S = (1+\epsilon)\Pr(A_s)$

    $x_T = e^{c_3 n^{\beta} \ln^2 n} \Pr(A_s)$, with $\beta = \frac{2}{k+1}$, $\epsilon > 0$
  - Apply LLL
- Best until 2010
  - Bohman&Keevash: $R(k,t) \geq c\left(\frac{t}{\ln t}\right)^{\frac{k+1}{2}}(\ln t)^{\frac{1}{k-2}}$

Major open problem

- Determine $\alpha(k)$ s.t. $R(k,t) = t^{\alpha(k)+o(1)}$
- Spencer 1975:$\alpha(k) \geq \frac{\binom{k}{2}-2}{k-2}$
- Spencer 1977:$\alpha(k) \geq \frac{k+1}{2} = \frac{\binom{k}{2}-1}{k-2}$
    - Best for 40+ years
    - How tight is it?
- $\alpha(k) \leq k-1$ since $R(k,t) \leq \binom{k+t-2}{k-1}$
- Conjecture: $\alpha(k)k-1$
    - Yes for $k=3$
    - Known for larger k

- This local lemma is so strong. Is it ultimate?



Local lemma is to determine a curve surrounding a *safe zone.*
Safe: $\Pr(\cap_i \overline{A_i}) > 0$ any set of events with the probabilities

Tight Bound of Lovász local lemma

- General (Non-symmetrical) case
- $\Pr(\cap \overline{A_i}) > 0$ if

$$\forall S \in \mathrm{ind}(G),\ \sum_{\substack{T \supseteq S \\ T \in \mathrm{ind}(G)}} (-1)^{|T \setminus S|} \prod_{i \in T} p_i > 0$$

- By James B. Shearer @IBM in 1985

James Shearer

Tight Bound of Lovász local lemma

- Symmetrical case
- $\Pr(\cap \overline{A_i}) > 0$ if

$$p < \begin{cases} \dfrac{(d-1)^{d-1}}{d^d} & \text{when } d > 1 \\ \dfrac{1}{2} & \text{when } d = 1 \end{cases}$$

- Corollary: $\Pr(\cap A_i) > 0$ if $edp \leq 1$

Application

- Any (k,s)-CNF is satisfiable if $s \leq \frac{1}{e}\frac{2^k}{k}$
  - Known: satisfiable if $s \leq \frac{1}{4}\frac{2^k}{k}$
  - Tight bound of s: $\left(\frac{2}{e} + o\left(\frac{1}{\sqrt{k}}\right)\right)\frac{2^k}{k}$ [Gebauer et al. 2011]
  - Can we efficiently find a satisfying assignment?

Algorithmic aspects

- **Like** other probabilistic methods, LLL proves existence non-constructively
- **Unlike** other probabilistic methods, LLL doesn't lead to efficient algorithms
  - Directly sampling has an *exponentially small* lower bound of success probability
  - Say, $\Pr(\cap \overline{A_i}) \geq \prod(1 - x_i)$ for general version
- Is there an efficient, constructive proof?

Constructive Lovász Local Lemma

- Initiated by Joszef Beck in 1991
  - Under strong conditions on neighborhood size
  - In terms of coloring, SAT …
- Breakthrough by Robin Moser&Gabor Tardos in 2009, Kashyap Kolipaka and Mario Szegedy in 2011
  - <u>Events are generated by independent random variables</u>
  - If Shearer's condition is met, an assignment s.t. none events occurs can be found in linear time

Joszef Beck

Gabor Tardos

Mario Szegedy

The assignment algorithm

For $X \in \mathcal{X}$ do

$\quad v_X \leftarrow$ a random evaluation of X

EndFor

While (some A occurs) do

$\quad$ Arbitrarily pick an event A that occurs

$\quad$ For $X \in \mathrm{vbl}(A)$ do

$\quad\quad v_X \leftarrow$ a random evaluation of X

$\quad$ EndFor

EndWhile

Return $(v_X)_{X \in \mathcal{X}}$

- $\mathrm{vbl}\ (A) \subset \Upsilon$: the set of variables determining A

Directions of LLL research

- Local conditions
  - Cluster LLL
  - Random walk
- Algorithms (Inspired by Moser&Tardos)
  - Efficient beyond Shearer's bound?
  - Efficient for abstract events?

Comparing probabilistic methods

- All dependent vs almost independent
  - Counting (union bound): mutually exclusive
  - First moment: linearity doesn't care dependence
  - Second moment: pairwise dependence
  - LLL: global dependence

References

- Spencer. Ramsey's theorem-A new lower bound. 1975
- Spencer. Asymptotic lower bounds for Ramsey functions. 1977
- James B. Shearer. On a Problem of Spencer. 1985
- Robin Moser and Gabor Tardos. A constructive proof of the general Lovasz Local Lemma. 2009
- Polipaka and Szegidy. Moser and Tardos Meet Lovász. 2011

- http://www.openproblemgarden.org/

Thank you

Appendix: Proof of Symmetric LLL

- Standard trick
  - Chain rule: $\Pr(\cap_i \overline{A_i}) = \prod_{i=1}^{n} \Pr(\overline{A_i} | \cap_{j=1}^{i-1} \overline{A_j})$
    - Valid only if each $\cap_{j=1}^{i-1} \overline{A_j}$ has nonzero probability
  - Hold if each term $\Pr\left(\overline{A_i} \middle| \cap_{j=1}^{i-1} \overline{A_j}\right) > 0$
- **Claim**: for any $t \geq 0$ and $A, B_1, B_2, \ldots B_t \in S$,

  1. $\Pr(\cap_{j=1}^{t} \bar{B}_j) > 0$

  2. $\Pr\left(A \middle| \cap_{j=1}^{t} \bar{B}_j\right) < x_i$

Inductive proof of the claim

- **Basis**: $t = 0$. Both 1 and 2 of the claim hold
- **Hypothesis**: the claim holds for all $t' < t$
- **Induction**
  - For **1**, $\Pr\left(\bigcap_{j=1}^{t} \overline{B_j}\right)$
    $$= \Pr\left(\overline{B_t} \mid \bigcap_{j=1}^{t-1} \overline{B_j}\right) \Pr\left(\bigcap_{j=1}^{t-1} \overline{B_j}\right) > 0$$
  - For **2**, let $\{C_1, \dots C_x\} = \{B_1, \dots B_t\} \cap \Gamma(A)$, and
    $$\{D_1, \dots D_y\} = \{B_1, \dots B_t\} \backslash \Gamma(A)$$
    - $x \leq d, x + y = t$

Proof: induction for 2

- If $x = 0$, A is independent of $\{B_1, \ldots B_t\}$ and $\Pr\left(A \middle| \cap_{j=1}^{t} \overline{B_j}\right) = \Pr(A) < \frac{1}{2d}$
- Assume $x > 0$. Then $y < t$.
  - $\Pr\left(A \middle| \cap_{j=1}^{t} \overline{B_j}\right) = \frac{\Pr(A \cap (\cap_{j=1}^{t} \overline{B_j}))}{\Pr(\cap_{j=1}^{t} \overline{B_j})}$

$$\leq \frac{\Pr\left(A \cap (\cap \overline{D_j})\right)}{\Pr\left((\cap \overline{C_j}) \cap (\cap \overline{D_j})\right)} = \frac{\Pr\left(A | \cap \overline{D_j}\right)}{\Pr\left((\cap \overline{C_j}) | \cap \overline{D_j}\right)}$$

Use chain rule to lower-bound $\Pr\left(\cap \overline{C_j} \mid \cap \overline{D_j}\right)$

General case